

The Language GFC

BNF-converter

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This document was automatically generated by the *BNF-Converter*. It was generated together with the lexer, the parser, and the abstract syntax module, which guarantees that the document matches with the implementation of the language (provided no hand-hacking has taken place).

The lexical structure of GFC

Identifiers

Identifiers $\langle Ident \rangle$ are unquoted strings beginning with a letter, followed by any combination of letters, digits, and the characters `_ ' ,` reserved words excluded.

Literals

String literals $\langle String \rangle$ have the form `"x"`, where x is any sequence of any characters except `"` unless preceded by `\`.

Integer literals $\langle Int \rangle$ are nonempty sequences of digits.

Double-precision float literals $\langle Double \rangle$ have the structure indicated by the regular expression $\langle digit \rangle + \cdot \langle digit \rangle + ('e' | 'E' | '-' | '+' \langle digit \rangle) ?$ i.e. two sequences of digits separated by a decimal point, optionally followed by an unsigned or negative exponent.

Reserved words and symbols

The set of reserved words is the set of terminals appearing in the grammar. Those reserved words that consist of non-letter characters are called symbols, and they are treated in a different way from those that are similar to identifiers. The lexer follows rules familiar from languages like Haskell, C, and Java, including longest match and spacing conventions.

The reserved words used in GFC are the following:

Ints	Str	Type
abstract	cat	concrete
data	flags	fun
grammar	in	lin
lincat	of	open
oper	param	pre
resource	table	transfer
variants		

The symbols used in GFC are the following:

```

;      =   {
}      :   ->
**     [   ]
\      .   (
)      -   <
>     $   ?
=>    !   ++
/      @   +
|      ,

```

Comments

There are no single-line comments in the grammar.

There are no multiple-line comments in the grammar.

The syntactic structure of GFC

Non-terminals are enclosed between \langle and \rangle . The symbols $::=$ (production), $|$ (union) and ϵ (empty rule) belong to the BNF notation. All other symbols are terminals.

```

<Canon> ::= grammar <ListIdent> of <Ident> ; <ListModule>
          |   <ListModule>

<Line>  ::= grammar <ListIdent> of <Ident> ;
          |   <ModType> = <Extend> <Open> {
          |   <Flag> ;
          |   <Def> ;
          |   }

<Module> ::= <ModType> = <Extend> <Open> { <ListFlag> <ListDef> }

```

```

⟨ModType⟩ ::= abstract ⟨Ident⟩
           | concrete ⟨Ident⟩ of ⟨Ident⟩
           | resource ⟨Ident⟩
           | transfer ⟨Ident⟩ : ⟨Ident⟩ -> ⟨Ident⟩

⟨ListModule⟩ ::= ε
              | ⟨Module⟩ ⟨ListModule⟩

⟨Extend⟩ ::= ⟨ListIdent⟩ **
           | ε

⟨Open⟩ ::= open ⟨ListIdent⟩ in
        | ε

⟨Flag⟩ ::= flags ⟨Ident⟩ = ⟨Ident⟩

⟨Def⟩ ::= cat ⟨Ident⟩ [ ⟨ListDecl⟩ ] = ⟨ListCIdent⟩
       | fun ⟨Ident⟩ : ⟨Exp⟩ = ⟨Exp⟩
       | transfer ⟨Ident⟩ = ⟨Exp⟩
       | param ⟨Ident⟩ = ⟨ListParDef⟩
       | oper ⟨Ident⟩ : ⟨CType⟩ = ⟨Term⟩
       | lincat ⟨Ident⟩ = ⟨CType⟩ = ⟨Term⟩ ; ⟨Term⟩
       | lin ⟨Ident⟩ : ⟨CIdent⟩ = \ ⟨ListArgVar⟩ -> ⟨Term⟩ ; ⟨Term⟩
       | ⟨Ident⟩ ⟨Status⟩ in ⟨Ident⟩

⟨ParDef⟩ ::= ⟨Ident⟩ ⟨ListCType⟩

⟨Status⟩ ::= data
          | ε

⟨CIdent⟩ ::= ⟨Ident⟩ . ⟨Ident⟩

⟨Exp1⟩ ::= ⟨Exp1⟩ ⟨Exp2⟩
        | ⟨Exp2⟩

⟨Exp⟩ ::= ( ⟨Ident⟩ : ⟨Exp⟩ ) -> ⟨Exp⟩
       | \ ⟨Ident⟩ -> ⟨Exp⟩
       | { ⟨ListEquation⟩ }
       | ⟨Exp1⟩

⟨Exp2⟩ ::= ⟨Atom⟩
        | data
        | ( ⟨Exp⟩ )

⟨Sort⟩ ::= Type

⟨Equation⟩ ::= ⟨ListAPatt⟩ -> ⟨Exp⟩

```

$$\begin{aligned}
\langle APatt \rangle & ::= (\langle CIdent \rangle \langle ListAPatt \rangle) \\
& | \langle Ident \rangle \\
& | \langle String \rangle \\
& | \langle Integer \rangle \\
& | \langle Double \rangle \\
& | - \\
\langle ListDecl \rangle & ::= \epsilon \\
& | \langle Decl \rangle \\
& | \langle Decl \rangle ; \langle ListDecl \rangle \\
\langle ListAPatt \rangle & ::= \epsilon \\
& | \langle APatt \rangle \langle ListAPatt \rangle \\
\langle ListEquation \rangle & ::= \epsilon \\
& | \langle Equation \rangle ; \langle ListEquation \rangle \\
\langle Atom \rangle & ::= \langle CIdent \rangle \\
& | < \langle CIdent \rangle > \\
& | \$ \langle Ident \rangle \\
& | ? \langle Integer \rangle \\
& | \langle String \rangle \\
& | \langle Integer \rangle \\
& | \langle Sort \rangle \\
\langle Decl \rangle & ::= \langle Ident \rangle : \langle Exp \rangle \\
\langle CType \rangle & ::= \{ \langle ListLabelling \rangle \} \\
& | (\langle CType \rangle => \langle CType \rangle) \\
& | \langle CIdent \rangle \\
& | **Str** \\
& | **Ints** \langle Integer \rangle \\
\langle Labelling \rangle & ::= \langle Label \rangle : \langle CType \rangle \\
\langle Term2 \rangle & ::= \langle ArgVar \rangle \\
& | \langle CIdent \rangle \\
& | < \langle CIdent \rangle \langle ListTerm2 \rangle > \\
& | \$ \langle Ident \rangle \\
& | \{ \langle ListAssign \rangle \} \\
& | \langle Integer \rangle \\
& | \langle Double \rangle \\
& | \langle Tokn \rangle \\
& | [] \\
& | (\langle Term \rangle)
\end{aligned}$$

$$\begin{aligned}
\langle \text{Term1} \rangle & ::= \langle \text{Term2} \rangle . \langle \text{Label} \rangle \\
& | \text{table } \langle \text{CType} \rangle \{ \langle \text{ListCase} \rangle \} \\
& | \text{table } \langle \text{CType} \rangle [\langle \text{ListTerm2} \rangle] \\
& | \langle \text{Term1} \rangle ! \langle \text{Term2} \rangle \\
& | \text{variants } \{ \langle \text{ListTerm2} \rangle \} \\
& | \langle \text{Term2} \rangle \\
\langle \text{Term} \rangle & ::= \langle \text{Term} \rangle ++ \langle \text{Term1} \rangle \\
& | \langle \text{Term1} \rangle \\
\langle \text{Tokn} \rangle & ::= \langle \text{String} \rangle \\
& | [\text{pre } \langle \text{ListString} \rangle \{ \langle \text{ListVariant} \rangle \}] \\
\langle \text{Assign} \rangle & ::= \langle \text{Label} \rangle = \langle \text{Term} \rangle \\
\langle \text{Case} \rangle & ::= \langle \text{ListPatt} \rangle => \langle \text{Term} \rangle \\
\langle \text{Variant} \rangle & ::= \langle \text{ListString} \rangle / \langle \text{ListString} \rangle \\
\langle \text{Label} \rangle & ::= \langle \text{Ident} \rangle \\
& | \$ \langle \text{Integer} \rangle \\
\langle \text{ArgVar} \rangle & ::= \langle \text{Ident} \rangle @ \langle \text{Integer} \rangle \\
& | \langle \text{Ident} \rangle + \langle \text{Integer} \rangle @ \langle \text{Integer} \rangle \\
\langle \text{Patt} \rangle & ::= (\langle \text{CIdent} \rangle \langle \text{ListPatt} \rangle) \\
& | \langle \text{Ident} \rangle \\
& | - \\
& | \{ \langle \text{ListPattAssign} \rangle \} \\
& | \langle \text{Integer} \rangle \\
& | \langle \text{Double} \rangle \\
\langle \text{PattAssign} \rangle & ::= \langle \text{Label} \rangle = \langle \text{Patt} \rangle \\
\langle \text{ListFlag} \rangle & ::= \epsilon \\
& | \langle \text{Flag} \rangle ; \langle \text{ListFlag} \rangle \\
\langle \text{ListDef} \rangle & ::= \epsilon \\
& | \langle \text{Def} \rangle ; \langle \text{ListDef} \rangle \\
\langle \text{ListParDef} \rangle & ::= \epsilon \\
& | \langle \text{ParDef} \rangle \\
& | \langle \text{ParDef} \rangle | \langle \text{ListParDef} \rangle \\
\langle \text{ListCType} \rangle & ::= \epsilon \\
& | \langle \text{CType} \rangle \langle \text{ListCType} \rangle
\end{aligned}$$

$$\begin{aligned}
\langle \text{ListCIdent} \rangle & ::= \epsilon \\
& \quad | \quad \langle \text{CIdent} \rangle \langle \text{ListCIdent} \rangle \\
\langle \text{ListAssign} \rangle & ::= \epsilon \\
& \quad | \quad \langle \text{Assign} \rangle \\
& \quad | \quad \langle \text{Assign} \rangle ; \langle \text{ListAssign} \rangle \\
\langle \text{ListArgVar} \rangle & ::= \epsilon \\
& \quad | \quad \langle \text{ArgVar} \rangle \\
& \quad | \quad \langle \text{ArgVar} \rangle , \langle \text{ListArgVar} \rangle \\
\langle \text{ListLabelling} \rangle & ::= \epsilon \\
& \quad | \quad \langle \text{Labelling} \rangle \\
& \quad | \quad \langle \text{Labelling} \rangle ; \langle \text{ListLabelling} \rangle \\
\langle \text{ListCase} \rangle & ::= \epsilon \\
& \quad | \quad \langle \text{Case} \rangle \\
& \quad | \quad \langle \text{Case} \rangle ; \langle \text{ListCase} \rangle \\
\langle \text{ListTerm2} \rangle & ::= \epsilon \\
& \quad | \quad \langle \text{Term2} \rangle \langle \text{ListTerm2} \rangle \\
\langle \text{ListString} \rangle & ::= \epsilon \\
& \quad | \quad \langle \text{String} \rangle \langle \text{ListString} \rangle \\
\langle \text{ListVariant} \rangle & ::= \epsilon \\
& \quad | \quad \langle \text{Variant} \rangle \\
& \quad | \quad \langle \text{Variant} \rangle ; \langle \text{ListVariant} \rangle \\
\langle \text{ListPattAssign} \rangle & ::= \epsilon \\
& \quad | \quad \langle \text{PattAssign} \rangle \\
& \quad | \quad \langle \text{PattAssign} \rangle ; \langle \text{ListPattAssign} \rangle \\
\langle \text{ListPatt} \rangle & ::= \epsilon \\
& \quad | \quad \langle \text{Patt} \rangle \langle \text{ListPatt} \rangle \\
\langle \text{ListIdent} \rangle & ::= \epsilon \\
& \quad | \quad \langle \text{Ident} \rangle \\
& \quad | \quad \langle \text{Ident} \rangle , \langle \text{ListIdent} \rangle
\end{aligned}$$